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distributions On approximating point spread

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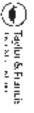
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On approximating point spread distributions

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fixed and small, and n is appropriately bounded, then the point spread distribution only slightly changes for small point differences. We also prove that for equal success rates p, the probability of a tie is minimized if p = 1/2. Numerical examples are included for the case with n = 12. We discuss some properties of the point spread distribution, defined as the distribution of the difference of two independent binomial random variables with the same parameter n including exact and approximate probabilities and related optimization issues. We use various approximation techniques for different distributions, special functions, and analytic, combinatorial and symbolic methods, such as multi-summation We prove that in case of unequal success rates, if these rates change with their difference kept

Keywords: Skellam distribution; approximating distributions; asymptotic enumeration; special functions;

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Introduction

beat the projected point spread through the bribing of players from the favoured team by gamblers. used to set fairly equal winning odds in terms of scoring differences in matches between players hockey, and soccer, independent Poisson random variables might be used. Point spreads are often modelled by independent binomial random variables, while in low-scoring sports, e.g. baseball, ing problems. In high-scoring sports, e.g. basketball, the underlying scoring distributions can be betting purposes [1]. On the other hand, there are some infamous cases involving point-shaving to or teams of widely different strengths. In fact, bookmakers set a point spread to even the game for Questions regarding the point spread distribution in certain sports present a rich variety of interest-

as large as $d \ge 0$ points in n games with respective success rates p and $p - \varepsilon$. Section 2 is devoted players in n 'games', respectively. We assume that X and Y are independent, binomially distributed We study the probability $f_0(n, p, \varepsilon)$ to the discussion of the symmetry of the point spread distribution $f_d(n, p, \varepsilon)$ (cf. Theorem 2.1). random variables and introduce the probability $f_d(n, p, \varepsilon)$ of a point spread X-Y which is at least In the simplified model, let X and Y represent the number of points scored by two teams or $f_{d+1}(n, p, \varepsilon)$, in Section 3, and briefly describe a random-walk-based approach for $f_1(n, p, \varepsilon)$ of a tie and that of an exact point spread d.

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Section 5. In Section 7, we ask some optimization questions such as: for a fixed $\varepsilon \ge 0$, what value of p will maximize $f_d(n, p, \varepsilon)$, a question settled in Theorem 7.1 for the case with $f_0(n, p, 0)$. are summarized in Theorems 2.1, 5.1, and 7.1. The last section is concerned with the distribution of the absolute point spread. The main results success rates are close to one. We extend some of the calculations to negative values of d in Sections 5 and 6 by a normal distribution in Theorem 5.1 and by a Skellam distribution if the the calculation of the distribution in Section 4. The point spread distribution is approximated in

the point spread distribution. For purely illustrative purposes, we consider an example which touches upon some aspects of

scoring experience of either player. We are interested in the following probabilities. free-throws. You can assume that the scoring luck is totally independent of the previous or current Example Players 1 and 2 are two basketball players. Player 1 makes 65% of his free-throws, while Player 2 is even better and makes 75%. They have a contest in which they each shoot 12

- What is the probability that this free-throw contest will end in a tie? What is the probability that Player 2 will win?
- What are the chances that, say, Player 2 scores two more than Player 1?
- (D) How about at least two more?

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0.1533 by sum (dbinom (0:12,12,0.65) * dbinom (0:12,12,0.75)), for (B) 0.6249the most economical way and conveniently in S-PLUS, at least in terms of the frugala similar fashion. four significant digits. Other values related to the point spread distribution can be derived in is obtained by sum(dbinom(1:12,12,0.75)*pbinom(0:11,12,0.65)), for (C) ity of the necessary code. The answers can be calculated easily with most statistical software, and perhaps in In fact, the answers are one-liners in S-PLUS: for (A) we get

approach, and discuss related optimization problems to help explain this observation. Calculations other words, if there are some external circumstances that are equally influencing both players to found, somewhat surprisingly, that the answers obtained changed only slightly in (B) and (D). In similar settings within a given range but with the same difference in shooting success rates and for n = 12 are included to illustrate numerical aspects and features of the different approaches. very small degree. We apply different approximation and summation techniques, a random walk play better or worse, then for small differences, the point spread distribution will change only to a We focus on questions (B) and (D), in particular. After calculating these values, we tested

Now we prove an interesting symmetry property.

player with success rate $p - \varepsilon$ with $\varepsilon \ge 0$. The function $f_d(n, p, \varepsilon)$, $\varepsilon \le p \le 1$, is symmetric success rate p accumulates at least $d \ge 0$ points more than the other (equally strong or 'weaker') THEOREM 2.1 Let $f_d(n, p, \varepsilon)$ denote the probability that, in n trials, the 'stronger' player with about $(1+\varepsilon)/2$ for every $n \ge$

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$$f_d(n, p, \varepsilon) = \sum_{k=d}^n \binom{n}{k} p^k (1-p)^{n-k} \sum_{j=0}^{k-d} \binom{n}{j} (p-\varepsilon)^j (1-p+\varepsilon)^{n-j}. \tag{1}$$
Let us assume that $\varepsilon \le p \le 1$. The claim is that $f_d(n, p, \varepsilon) = f_d(n, 1-p+\varepsilon, \varepsilon)$. By Equation

$$f_d(n, 1-p+\varepsilon, \varepsilon) = \sum_{k=d}^n \binom{n}{k} (1-p+\varepsilon)^k (p-\varepsilon)^{n-k} \sum_{j=0}^{k-d} \binom{n}{j} (1-p)^j p^{n-j}.$$

We can view the second summation as the probability that the player with success rate p loses at most k-d games, while $\binom{n}{k}(1-p+\varepsilon)^k(p-\varepsilon)^{n-k}$ is the probability that the player with success rate $p-\varepsilon$ loses exactly k games. Thus, the combined expression gives the probability of winning by at least d games by the former player.

Note that this property guarantees that it suffices to deal with $f_d(n, p, \varepsilon)$ with $p \ge (1 + \varepsilon)/2$. Sometimes, we use the short notation f_d instead of $f_d(n, p, \varepsilon)$.

statement holds true for any entire function (whose Taylor series converges to the function itself of the polynomial f_d about $(1+\varepsilon)/2$ which has only terms with even exponents. Similarly, the the extended function remains symmetric. This can be proved by taking the Taylor expansion The definition of f_d can be extended to every real p by Equation (1), and obviously,

a polynomial in p, is $(-1)^d c_{n,d}$ with Remark 2 Note that the degree of polynomial f_d in p is 2n. Moreover, Equation (1) shows that $p^d(1-p+\varepsilon)^d$ is always a factor of f_d independent of n. By focusing on the highest powers of p (including sign) in Equation (1), it can be proved that the leading coefficient of $f_d(n, p, \varepsilon)$, as

$$c_{n,d} = {2n-1 \choose n-d}, \quad n \ge d,$$

independent of ε . It follows, for instance, that $c_{n,0} = c_{n,1} = {2n \choose n}/2$.

For d = 0, by more detailed calculations, we can get that

$$f_0(n, p, \varepsilon) = (1 + \varepsilon)^n - n(1 + \varepsilon)^{n-1} (1 + n\varepsilon) p$$

$$+ \frac{n}{4} (1 + \varepsilon)^{n-2} (\varepsilon^2 n^3 + 6\varepsilon n^2 - \varepsilon^2 n - 2\varepsilon n + 6n - 2) p^2 + \dots + \binom{2n-1}{n} p^{2n}.$$

rather than about p = 0. However, as we will see, it is more important to have the expansion of f_d about $p = (1 + \varepsilon)/2$

The probability of a tie and exact point spread

If p=1 then the probability of a tie is $(1-\varepsilon)^n$. In general, we find that the probability of a tie is

$$f = f_0(n, p, \varepsilon) - f_1(n, p, \varepsilon) = \sum_{k=0}^{n} {n \choose k} p^k (1-p)^{n-k} {n \choose k} (p-\varepsilon)^k (1-p+\varepsilon)^{n-k}$$

$$= \sum_{k=0}^{n} {n \choose k}^2 (p(p-\varepsilon))^k ((1-p)(1-p+\varepsilon))^{n-k}$$

$$= ((1-p)(1-p+\varepsilon))^n \sum_{k=0}^{n} {n \choose k}^2 \left(\frac{p(p-\varepsilon)}{(1-p)(1-p+\varepsilon)} \right)^k,$$

and thus,

$$T = \begin{cases} ((1-p)(1-p+\varepsilon))^n \left(1 - \frac{p(p-\varepsilon)}{(1-p)(1-p+\varepsilon)}\right)^n \\ \times P_n \left(\frac{1+p(p-\varepsilon)/(1-p)(1-p+\varepsilon)}{1-p(p-\varepsilon)/(1-p)(1-p+\varepsilon)}\right), & \text{if } \frac{1+\varepsilon}{2}$$

with the *n*th Legendre polynomial $P_n(x)$ (cf. [2]). If x < -1 then $P_n(x)$ can be approximated by

$$P_n(x) \sim \frac{(-1)^n}{\sqrt{2\pi n}(x^2 - 1)^{1/4}} (-x + \sqrt{x^2 - 1})^{n+1/2}$$

as $n \to \infty$. For instance, the answer 0.1533 to question (A) can be approximated by 0.1542 this way since p < 1 guarantees that the argument of P_n in (3) is less than -1. With the notation

$$\Delta = p - \frac{1+\varepsilon}{2},$$

we get the approximation of the answer T

$$T pprox \left(\Delta - rac{1}{2}
ight)^{2n} \left(1 - rac{(\Delta + 1/2)^2}{(\Delta - 1/2)^2}
ight)^n P_n \left(-rac{1}{4\Delta} - \Delta
ight)$$

for a small ε if $\Delta > 0$. Of course, we can also use the approximation (3) for the last factor. For example, with n=12, p=0.75, and $\varepsilon=0.01$, we get that the exact value is 0.1859 while the above approximation with Equation (3) results in 0.1857. Also note that by Equation (2), we have

$$f_0\left(n, \frac{1+\varepsilon}{2}, \varepsilon\right) - f_1\left(n, \frac{1+\varepsilon}{2}, \varepsilon\right) \sim \frac{(1-\varepsilon^2)^n}{\sqrt{n\pi}}$$

for $\Delta = 0$ and $n \to \infty$.

We can generalize the above approach for any difference $d \ge 0$. By using the hypergeometric function $_2F_1$ (cf. [2]), we obtain the following theorem.

THEOREM 3.1 The probability $f_d(n, p, \varepsilon) - f_{d+1}(n, p, \varepsilon)$ of an exact point spread of d is

$$\left(\left(1-p\right)\left(1-p+\varepsilon\right)\right)^{n} \left(\frac{p}{1-p}\right)^{d} \binom{n}{d} {}_{2}F_{1}\left(d-n,-n,1+d;\frac{p\left(p-\varepsilon\right)}{\left(1-p\right)\left(1-p+\varepsilon\right)}\right)$$

 $if (1+\varepsilon)/2 \le p < 1.$

For instance, Theorem 3.1 gives the answer $f_0(12, 0.75, 0.10) - f_1(12, 0.75, 0.10) = 0.1533$ to question (A). Similarly, $f_2(12, 0.75, 0.10) - f_3(12, 0.75, 0.10) = 0.1676$ answers ques-

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4. Random walk approach

Assume that in the shooting competition the players shoot in an alternating fashion. Let X_i be 1 or 0 if the stronger player scores or misses his free throw in the *i*th trial. We define Y_i similarly. Then $X = \sum_{k=1}^{n} X_i$ and $Y = \sum_{k=1}^{n} Y_i$. We can model the problem with a random walk which has step sizes $X_i - Y_i$, i = 1, 2, ..., n: the walk moves one step to the right or left if only the stronger or the other player scores, respectively, and stays in place if either both or none of the players score. With P = p(1-q), Q = q(1-p), and $q = p - \varepsilon$, we have that

$$P(X - Y = k) = [z^k] \left(Pz + (1 - P - Q) + \frac{Q}{z} \right)^n = [z^{n+k}] (Pz^2 + (1 - P - Q)z + Q)^n$$

where $[z^k]g(z)$ stands for the coefficient of the term z^k in the (Laurent) power series expansion Formally, in terms of a Laurent series with an essential singularity at 0, we can write that of g(z) about 0. By adding the coefficients for $k = d, d + 1, \ldots, n$, we can obtain $f_d(n, p, \varepsilon)$.

$$f_d(n, p, \varepsilon) = P^n \left[z^{n+d} \right] \left(1 + \frac{1}{z} + \frac{1}{z^2} + \cdots \right) \left(z^2 + \frac{1 - P - Q}{P} z + \frac{Q}{P} \right)^n$$

extraction is easy when p=q=1/2 and thus, $\varepsilon=0,\,P=Q=$ for extracting coefficients. In some cases though, it might be helpful. For example, [3, Proposition VII. 10] suggests immediately. In another example, we can get a good approximation for the probability of or no improvement over Equation (1) unless some computer algebra system is easily available T = P(X - Y = 0) (cf. Equation (2)) by using standard facts regarding random walks. In fact, We note, however, that this approach results in a double sum for f_d and thus, offers little 1/4, and identity (15) follows

$$T \sim \frac{\left(1 - (\sqrt{P} - \sqrt{Q})^2\right)^{n+1/2}}{2(PQ)^{1/4}\sqrt{\pi n}}$$

of the approximated T as n grows: it is exponential if $p \neq q$ and of order $1/\sqrt{n}$ if p = q. In way we get the approximate answer 0.1542 to (A). Note the difference in the rate of decrease approximation for T which looks simpler than the one given in Section 3. For example, in this probability that a random walk of length n is a bridge from altitude 0 to 0), hence providing an i.e. $f_0(n, p, \varepsilon)$ – to Wagon [5]. The author conjectures that for $n \ge 5$, $p = (1 + \varepsilon)/2$ switches from being the this particular location gives either a (local) minimum or maximum. location of the minimum of T to that of the maximum as ε grows. Of course, the symmetry of T, T decreases as $n \to \infty$, as can be shown in the way Lemma 5 of [4] is proved according $\rightarrow \infty$ by using the asymptotic number of 'bridges' (and its generalization for finding the $f_1(n, p, \varepsilon)$, in p about $(1 + \varepsilon)/2$ follows by Theorem 2.1, and this yields that

Approximation by normal

the range $p \ge (1+\varepsilon)/2$. Let X and Y represent the number of successful shots made by Players 2 and 1, respec- $\varepsilon > 0$ for Players 2 and 1, respectively, we first approximate $f_d(n, p, \varepsilon)$ for $d \ge 1$ and then extend the range of the approximation for $d \le 0$. As we observed above, we can restrict our attention to With some reasonable bounds on the number of trials n and success rates p and q =

tively. Clearly, X and Y are independent, binomially distributed random variables,

values of n, as long as $9p/(1-p) \le n$, i.e. $p \le n/(n+9)$, although this condition can be relaxed in practice. Therefore, the distribution of X-Y is approximately $\mathcal{N}[\mu=n(p-q)=n\varepsilon,\sigma=n]$ $\sqrt{n(p(1-p)+q(1-q))}$]. $\sqrt{np(1-p)}$] and $Y \sim \mathcal{N}[\mu = nq, \sigma = \sqrt{nq(1-q)}]$ approximately, even for small

To answer (A)–(D) we need the probabilities P(X-Y=0), $P(X-Y\ge1)$, P(X-Y=2), and $P(X-Y\ge2)$, respectively. For example, $f_1(n,p,\varepsilon)=P(X-Y\ge1)$ can be approximated by

$$1 - \Phi\left(\frac{0.5 - n\varepsilon}{\sqrt{n(p(1-p) + q(1-q))}}\right)$$

using the so-called continuity correction. We can use a simple and fairly accurate approximation for the normal distribution function

$$\Phi(x) \approx 0.5 + 0.1x(4.4 - x)$$
, for $0 \le x \le 2.2$

In fact, it is good to two decimal places.

We consider $\sqrt{n(p(1-p)+q(1-q))}$ from Equation (5). We first take p(1-p)+q(1-q) and rewrite it in terms of the fixed $\varepsilon=p-q$ and then express it as a function of $\Delta=p-(1+q)$

$$\sqrt{n(p(1-p)+q(1-q))} = \sqrt{\frac{n}{2}\left(1-4\left(p-\frac{1+\varepsilon}{2}\right)^2 - \varepsilon^2\right)}$$
$$\approx \sqrt{\frac{n}{2}\left(1-2\left(p-\frac{1+\varepsilon}{2}\right)^2 - \frac{1}{2}\varepsilon^2\right)}$$

if ε and Δ are small, more precisely, if $\sqrt{n}(\Delta^4 + \varepsilon^4)$ is small. We can use a quadratic or finer approximation of the argument of Φ about $p = (1 + \varepsilon)/2$ in Equation (5). We proceed with the quadratic approximation. Let us assume that $2\varepsilon \Delta^4 \sqrt{2n}$ is small. This assumption will guarantee that the above approximation introduces only negligible errors when we replace x by x_1, x_2 , and x_d below. We set

$$x_1 = \left(0.5\sqrt{\frac{2}{n}} - \varepsilon\sqrt{2n}\right)\left(1 + 2\left(p - \frac{1 + \varepsilon}{2}\right)^2 + \frac{\varepsilon^2}{2}\right) \le 0,$$

for $0.5/\varepsilon \le n \le 2/\varepsilon^2$ and use

$$1 - \Phi(x) \approx 0.5 + 0.1(-x)(4.4 + x) \tag{6}$$

with $x = x_1$. This provides us with a quadratic approximation in x_1 . Note that to answer (D), we need a slight modification of x_1 . We use approximation (6) with setting x to

$$x_2 = \left(1.5\sqrt{\frac{2}{n}} - \varepsilon\sqrt{2n}\right) \left(1 + 2\left(p - \frac{1+\varepsilon}{2}\right)^2 + \frac{\varepsilon^2}{2}\right) \le 0$$

for $1.5/\varepsilon \le n \le 2/\varepsilon^2$. In general, for a difference of d points we se

$$x_d = \left((d - 0.5) \sqrt{\frac{2}{n}} - \varepsilon \sqrt{2n} \right) \left(1 + 2 \left(p - \frac{1 + \varepsilon}{2} \right)^2 + \frac{\varepsilon^2}{2} \right) \le 0 \tag{7}$$

for $(d - 0.5)/\varepsilon \le n \le 2/\varepsilon^2$.

d = 1 we get that We use the above quadratic approximation (6) to $f_d(n, p, \varepsilon)$ about $(1 + \varepsilon)/2$. For example, for

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the constant coeff:
$$\left(0.5 - \frac{0.311127}{\sqrt{n}} - \frac{0.05}{n}\right) + \varepsilon \left(0.622254\sqrt{n} + 0.2\right)$$

 $-\varepsilon^2 \left(0.2n + \frac{0.155563}{\sqrt{n}} + \frac{0.05}{n}\right) + O\left(\varepsilon^3\right),$ (8)
the coeff of $\left(p - \frac{1+\varepsilon}{2}\right)^2$: $-\left(\frac{0.622254}{\sqrt{n}} + \frac{0.2}{n}\right) + \varepsilon \left(1.24451\sqrt{n} + 0.8\right)$
 $-\varepsilon^2 \left(0.8n + \frac{0.1}{n}\right) + O\left(\varepsilon^3\right).$ (9)

Note that if $d \le 0$ then $x_d < 0$ in Equation (7), and we need only the condition that $n \le 2/\varepsilon^2$. In general, we have the following theorem.

THEOREM 5.1 For an arbitrary integer d, we get the approximation $f_d(n, p, \varepsilon) \approx c_0 + c_2(p - 1)$

$$c_{0} = \left(0.5 - \frac{0.311127(2d - 1)}{\sqrt{n}} - \frac{0.05(4d(d - 1) + 1)}{n}\right) + \varepsilon(0.622254\sqrt{n} + 0.2(2d - 1))$$

$$- \varepsilon^{2} \left(0.2n + \frac{0.155563(2d - 1)}{\sqrt{n}} + \frac{0.05(4d(d - 1) + 1)}{n}\right) + O(\varepsilon^{3}), \qquad (10)$$

$$c_{2} = -\left(\frac{0.622254(2d - 1)}{\sqrt{n}} + \frac{0.2(4d(d - 1) + 1)}{n}\right) + \varepsilon(1.24451\sqrt{n} + 0.8(2d - 1))$$

$$- \varepsilon^{2} \left(0.8n + \frac{(2d - 1)^{2}}{10n}\right) + O(\varepsilon^{3}), \qquad (11)$$

provided that $(d-0.5)/\varepsilon \le n \le 2/\varepsilon^2$ and $2\varepsilon \Delta^4 \sqrt{2n}$ is small with $\Delta = p - (1+\varepsilon)/2$.

Of course, if p is close to $(1+\varepsilon)/2$ and ε is small then the constant term $0.5-0.311127(2d-1)/\sqrt{n}-0.05(4d(d-1)+1)/n$ suffices in order to get a good approximation

coefficients. Thus, it will not help in finding a better match for f_d . approximating function (7) since the Maclaurin series of $1/\sqrt{a-y}$, a>0, in y has only positive Note that using a higher degree approximation in Equation (5) will not change the shape of the

The last approximation (11) explains the changes in f_d , mentioned in Section 1, if we change

(11) after dropping the terms with $O(\varepsilon^3)$. Note that in some cases the condition $(d-0.5)/\varepsilon \le n$ is violated and thus, the approximation fails to reach an acceptable accuracy. p but keep n, ε , and d fixed.

Tables 1 and 2 show the exact probabilities and their approximations via Equations (10) and

Poisson approximation

for large n. In this case, one might try to use approximation by the Poisson distribution for n-X and n-Y, respectively. For instance, if $p-\varepsilon \ge 0.90$ and $n(1-p+\varepsilon) \le 10$ then If $p - \varepsilon$ is close to 1 then the approximation (5) by normal distribution does not work except $p - \varepsilon \ge 0.90$ and $n(1 - p + \varepsilon) \le 10$ then

Table 1. The values of $f_1(12, p, \varepsilon)$ with $p = 0.55, 0.60, \dots, 0.80$ and $\varepsilon = 0.01, 0.05, 0.10$.

		$\varepsilon = 0.01$		$\varepsilon = 0.05$		$\varepsilon = 0.10$
p	Exact	Approximation	Exact	Approximation	Exact	Approximation
0.55	0.4386	0.4290	0.5176	0.5177	0.6151	0.6171
0.60	0.4376	0.4280	0.5178	0.5179	0.6156	0.6176
0.65	0.4359	0.4263	0.5180	0.5182	0.6173	0.6193
0.70	0.4332	0.4238	0.5185	0.5188	0.6203	0.6220
0.75	0.4291	0.4206	0.5190	0.5195	0.6249	0.6258
0.80	0.4228	0.4166	0.5198	0.5203	0.6317	0.6308

2. The values of $f_2(12, p, \varepsilon)$ with p = 0.55, 0.60, ..., 0.80 and $\varepsilon = 0.01, 0.05, 0.10$.

	~	$\varepsilon = 0.01$	~	$\varepsilon = 0.05$	~	$\varepsilon = 0.10$
р	Exact	Approximation	Exact	Approximation	Exact	Approximation
0.55	0.2864	0.2191	0.3579	0.3242	0.4539	0.4429
0.60	0.2835	0.2147	0.3563	0.3223	0.4536	0.4426
0.65	0.2782	0.2073	0.3531	0.3185	0.4527	0.4417
0.70	0.2699	0.1967	0.3479	0.3128	0.4509	0.4403
0.75	0.2578	0.1830	0.3401	0.3053	0.4482	0.4383
0.80	0.2400	0.1662	0.3283	0.2959	0.4439	0.4357

the distribution of X-Y can be approximated by the distribution of the difference of two indepenfollows a Skellam distribution. We have that dent Poisson random variables n $X \sim \text{Poisson}[n(1-p)]$ and $n-Y \sim \text{Poisson}[n(1-p+\varepsilon)]$, approximately. Therefore Y and n - X. The difference of two Poisson random variables

$$f_d(n, p, \varepsilon) \approx P(X - Y \ge d) = P((n - Y) - (n - X) \ge d)$$

$$= \sum_{k \ge d} e^{-(\mu_1 + \mu_2)} \left(\frac{\mu_1}{\mu_2}\right)^{k/2} I_k(2\sqrt{\mu_1 \mu_2})$$

yields 0.2283. with $\mu_1 = n(1 - p + \varepsilon)$ and $\mu_2 = n(1 - p)$, and $I_k(x)$ being the modified Bessel function of the first kind [2]. For example, we get that $f_2(12, 0.95, 0.05) = 0.2247$ while the above approximation

7. Optimization

kept fixed. Clearly, One might wonder what the largest possible probabilities in (B) and (D) are. Note that using the probabilistic context, $f_d(n, p, \varepsilon)$ increases as $n \to \infty$ and $\varepsilon > 0$, $d \ge 0$, and $p, \varepsilon \le p \le 1$ are

$$f_0(n, p, 0) = 1 - f_1(n, p, 0),$$
 (12)

and its value, (1+T)/2, can be determined by identity (2), possibly using the approximation indefinitely. (3) for T. By Remark 2, $f_d(n, p, \varepsilon)$ goes to ∞ for d even and to $-\infty$ for d odd as |p| grows

It is more interesting to look for

$$\max_{\varepsilon \le p \le 1} f_d(n, p, \varepsilon)$$

for different values of d with n and ε kept fixed. Of course, by symmetry, we can, as we will do, focus on the range $(1 + \varepsilon)/2 \le p \le 1$.

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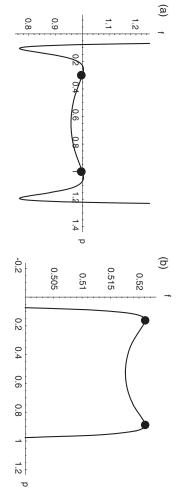


Figure 1. The largest probabilities in two examples with the maximum locations emphasized

For other values of ε this change at $p=(1+\varepsilon)/2$ happens at a higher value of d depending on n and ε , too. For example, with $\varepsilon=0.3$, $(\partial^2/\partial p^2)f_d(12,0.65,0.3)$, $(\partial^2/\partial p^2)f_d(13,0.65,0.3)$, and $(\partial^2/\partial p^2)f_d(16,0.65,0.3)$ become negative at d=4,5, and 6, respectively. fairly flat about its vertex for all $n \ge 1$. The shape becomes concave down at p = 1/2 for $d \ge 1$. down as d increases. For $\varepsilon = 0$, the shape of f_0 is that of a distorted parabola opening up which is We observe that the shape of f_d about $p = (1 + \varepsilon)/2$ changes from concave up to concave

at $p = (1 + \varepsilon)/2$. Figure 1(b), $f_1(12, p, 0.05)$). On the other hand, as d grows, it appears that the maximum occurs 0.8883, and the probability appears to be sharply decreasing as p increases from this value on (cf. 6 can hardly help. In fact, for n = 12, $\varepsilon = 0.05$, and d = 1 the optimum is found around p = 0.05As a consequence, typically, the maximum occurs at p=1 when d is small (cf. Figure 1(a), $f_0(12, p, 0.3)$). However, when it is not the case, the approximation methods of Sections 5 and

We prove only the following theorem.

THEOREM 7.1 The polynomial $f_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and thus, it takes its minimum at $p_0(n, p, 0)$ is concave up, and

Proof By the notation (4), we get that

$$f_d(n, p, \varepsilon) = \left(\left(\frac{1 - \varepsilon}{2} - \Delta \right) \left(\frac{1 + \varepsilon}{2} - \Delta \right) \right)^n \sum_{k=d}^{\infty} \binom{n}{k} \left(\frac{\Delta + (1 + \varepsilon)/2}{(1 - \varepsilon)/2 - \Delta} \right)^n$$
$$\times \sum_{j=0}^{k-d} \binom{n}{j} \left(\frac{\Delta + (1 - \varepsilon)/2}{(1 + \varepsilon)/2 - \Delta} \right)^j.$$

With $\varepsilon = 0$, this simplifies to

$$\frac{(1-2\Delta)^{2n}}{2^{2n}} \sum_{k=d}^{n} \binom{n}{k} \left(\frac{2\Delta+1}{1-2\Delta}\right)^{k} \sum_{j=0}^{k-d} \binom{n}{j} \left(\frac{2\Delta+1}{1-2\Delta}\right)^{j},\tag{13}$$

which can be expanded as a function of Δ^2

$$f_d(n, p, 0) = \sum_{k=0}^{n} c_{2k}(n, d) \Delta^{2k}$$

according to Theorem 2.1. In general, determining $c_{2k}(n,d)$ requires the evaluation of a triple sum according to Formula (13). On the other hand, we can easily derive that the coefficient of the

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term $\Delta^2 = (p - 1/2)^2$ is

$$c_2(n,d) = \frac{1}{2^{2n-2}} \sum_{k=d}^n \sum_{j=0}^{k-d} \binom{n}{k} \binom{n}{j} (2(k+j-n)^2 - n).$$

For instance, $c_2(n, n) = -n/2^{2n-2} < 0$ follows immediately. With some calculations and simplifications, we obtain the recurrence relation

$$c_2(n, d+1) = c_2(n, d) - \frac{1}{2^{2n-2}} \sum_{k=d}^n \binom{n}{k} \binom{n}{k-d} (2(2k-d-n)^2 - n), \text{ for } d \ge n$$

with the initial condition

$$c_2(n,0) = \frac{1}{2^{2n-2}} \binom{2n-2}{n-1} \sim \frac{1}{\sqrt{\pi(n-1)}},\tag{1}$$

which already guarantees the convexity of $f_0(n, p, 0)$ at p = 1/2. We note that Equation (14) by Equations (2) and (12), $c_0(n, 1) = 1 - c_0(n, 0)$ implies that follows from the observation (12) which implies that $c_2(n, 1) = -c_2(n, 0)$. In a similar fashion,

$$f_0\left(n, \frac{1}{2}, 0\right) = c_0(n, 0) = \frac{1}{2}\left(1 + \frac{\binom{2n}{n}}{2^{2n}}\right).$$

With considerably more calculations, we obtain all coefficients

$$c_{2k}(n,0) = \frac{1}{2^{2n-2k}} \binom{2k-1}{k} \binom{2n-2k}{n-k}, \quad k = 1, 2, \dots, n,$$
 (16)

p = 1/2, while the locations of the maximum are 0 and 1 since the function is symmetric about which proves the convexity of $f_0(n, p, 0)$ everywhere. The minimum of $f_0(n, p, 0)$ is taken at p = 1/2 by Theorem 2.1.

Note that Equation (16) can be verified by finding that

$$(2(n+2)N - (2n+1))(2(n+2)N^2 - (2n+3)(1+4\Delta^2)N + 8(n+1)\Delta^2)$$

package, for the Zb or the certificate finding FindRecurrence function of the MultiSum [7] Mathematical contents of MultiSum [7] Mathema is an annihilating operator (cf. [6,7]), using Zeilberger's algorithm [6] by calling the Zb function of

$$\sum_{k=1}^{n} \frac{1}{2^{2n-2k}} \binom{2k-1}{k} \binom{2n-2k}{n-k} \Delta^{2k} = \frac{\binom{2n}{2}}{2^{2n+1}} \left(-1 + {}_{2}F_{1}\left(\frac{1}{2}, -n, \frac{1}{2} - n, (2\Delta)^{2}\right)\right), \quad (17)$$

Mathematica package Sigma as it was pointed out by Schneider [8]. The proof is complete after the right factor annihilates the sums up to an inhomogeneous part which is free of n; in fact, it comparing the initial values of Equation (17) and the double sum $f_0(n, p, 0) - c_0(n, 0)$ based on left factor of the annihilator changes to N-1 if we include the constant term, indicating that with N being the forward shift operator with respect to n. It also annihilates the double sum $f_0(n, p, 0) - c_0(n, 0)$ which can be numerically verified for particular values of n. (Note that the $4x^2$)/2.) The symbolic verification can be effectively done by the latest version of the

We note that for d = 0, Formula (13) simplifies to

$$f_0(n, p, 0) = \frac{(1 - 2\Delta)^{2n}}{2^{2n}} \sum_{k=0}^{n} \sum_{j=0}^{n} \binom{n}{k} \binom{n}{k+j} \left(\frac{2\Delta + 1}{1 - 2\Delta}\right)^{2k+j}.$$

becomes concave down at p = 1/2. We add that Equation (12) also implies that $c_{2k}(n, 1) = -c_{2k}(n, 0), k \ge 1$, and thus, $f_1(n, p, 0)$

Note that for d = n it is straightforward to prove the following theorem.

1, 1, and n - 1, respectively. Thus,The roots of $(\partial/\partial p) f_n(n, p, \varepsilon)$ are 0, $(1+\varepsilon)/2$, and $1+\varepsilon$ with multiplicity n

$$\max_{\varepsilon \le p \le 1} f_n(n, p, \varepsilon) = f_n\left(n, \frac{1+\varepsilon}{2}, \varepsilon\right)$$

We note that very recently the problem of finding

$$\max_{n}(1-f_0(n,p,\varepsilon))$$

that the weaker player wins by more than $d \ge 0$ points and the corresponding n are close to $0.5 - 0.88\sqrt{\varepsilon(1+2d)}$ and $(1+2d)/(2\varepsilon)$, respectively, for any sufficiently small $\varepsilon > 0$ and $|\Delta|$. the weaker player, with conveniently switching the sign of d, we get the maximum probability The results can be easily generalized for $1 - f_d(n, p, \varepsilon)$ with the difference $d \le 0$, i.e. the probability that the weaker player scores at least $-d + 1 \ge 1$ points more than the stronger one, since and $0.5-0.88\sqrt{\varepsilon}$, respectively, for any sufficiently small $\varepsilon>0$ and $|\Delta|$. We observe that only one. The constant term approximation (10) of Theorem 5.1 and numerical evidence suggest that the optimum value n and the corresponding probability $1 - f_0(n, p, \varepsilon)$ must be close to $1/(2\varepsilon)$ in this case the approximation (7) works even for small values of n. From the point of view of ε seems to matter as long as $|\Delta|$ is small, in agreement with our findings regarding (B) and (D). $f_0(n, p, \varepsilon)$ corresponds to the probability that the weaker player scores more than the stronger appeared in [9,4] with the pairs p = 0.51, $\varepsilon = 0.01$, and p = 0.101, $\varepsilon = 0.001$. Here 1 –

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maximization problem open. from the location of the maximum to that of the minimum as ε grows, leaving the original $(1+\varepsilon)/2$. In general, for $f_d(n, p, \varepsilon)$ with $d \ge 1$, the author believes that $p = (1+\varepsilon)/2$ switches On another note, Wagon [5] conjectures that $f_0(n, p, \varepsilon), \varepsilon \le p \le 1$, takes its minimum at

3. The absolute spread difference for p = q = 1/2

We can consider the absolute spread difference between X and Y. Here we deal with the special case p=q=1/2 which implies $\varepsilon=0$. Since $X-Y\sim \mathcal{N}[\mu=0,\sigma=\sqrt{n}/2]$ approximately, we get that |X-Y| is approximately of half-normal distribution, i.e. the distribution of the absolute moments of |X - Y| raises some interesting questions involving double summations value of a normally distributed random variable centred at zero with $\sigma = \sqrt{n/2}$. Determining the

THEOREM 8.1 For the rth raw moment of |X - Y| we get that

$$m_r = \sum_{k=0}^n \sum_{j=0}^n \frac{|k-j|^r}{2^{2n}} \binom{n}{k} \binom{n}{j} = 2 \sum_{k=0}^n \sum_{j=0}^k \frac{(k-j)^r}{2^{2n}} \binom{n}{k} \binom{n}{j}$$
(18)

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with the particular values

$$m_0 = 1 - \frac{\binom{2n}{n}}{2^{2n}}, \quad m_1 = \frac{n}{2^{2n}} \binom{2n}{n}, \quad m_2 = \frac{n}{2}, \quad m_3 = \frac{n^2}{2^{2n}} \binom{2n}{n},$$
 $m_4 = \frac{n(3n-1)}{4}, \quad m_5 = \frac{n^2(2n-1)}{2^{2n}} \binom{2n}{n}, \quad \text{and} \quad m_6 = \frac{n(15n^2 - 15n + 4)}{8}.$

Of course, $E(|X - Y|^2) = \text{var}(X - Y) = n/2$. We also observe that $m_0 = 1 - P(X = Y) = 1 - \binom{2n}{n}/2^{2n}$ which provides us with an alternative proof of Equation (15). The proof of Theorem 8.1 can be accomplished by the use of the Mathematica package MultiSum (cf. [7]) or Sigma (cf. [10]). Further values of m_r have been determined by Schneider second form. The closed form for m_1 was originally suggested by John Essam and derived in [7]. have standard bounds if we use the first form of m_r in Equation (18) while odd indices require the [8]. Note that finding moments with even indices might be easier since the summation variables

distribution in an asymptotic sense. We note that Theorem 8.1 is in agreement with the moments of the corresponding half-normal

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